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LERNER, DAVID, LITTENBERG, KRUMHOLZ & MENTLIK 600 SOUTH AVENUE WEST WESTFIELD, NJ 07090			SHAW, YIN CHEN	
			ART UNIT	PAPER NUMBER
			2135	

DATE MAILED: 11/02/2006

Please find below and/or attached an Office communication concerning this application or proceeding.

<b>Office Action Summary</b>	<b>Application No.</b> 10/069,176	<b>Applicant(s)</b> ASANO ET AL.	
	<b>Examiner</b> Yin-Chen Shaw	<b>Art Unit</b> 2135	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

**Period for Reply**

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

**Status**

- 1) ☒ Responsive to communication(s) filed on 07 August 2006.
- 2a) ☒ This action is **FINAL**.                      2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

**Disposition of Claims**

- 4) ☒ Claim(s) 34-73 is/are pending in the application.
- 4a) Of the above claim(s) \_\_\_\_\_ is/are withdrawn from consideration.
- 5) ☐ Claim(s) \_\_\_\_\_ is/are allowed.
- 6) ☒ Claim(s) 34-73 is/are rejected.
- 7) ☐ Claim(s) \_\_\_\_\_ is/are objected to.
- 8) ☐ Claim(s) \_\_\_\_\_ are subject to restriction and/or election requirement.

**Application Papers**

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on \_\_\_\_\_ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.  
 Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).  
 Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

**Priority under 35 U.S.C. § 119**

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All    b) ☐ Some \* c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
  2. ☐ Certified copies of the priority documents have been received in Application No. \_\_\_\_\_.
  3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

\* See the attached detailed Office action for a list of the certified copies not received.

**Attachment(s)**

- |  |   |
|--|---|
| 1) <input checked="" type="checkbox"/> Notice of References Cited (PTO-892)                                | 4) <input type="checkbox"/> Interview Summary (PTO-413)<br>Paper No(s)/Mail Date. _____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948)                       | 5) <input type="checkbox"/> Notice of Informal Patent Application                       |
| 3) <input type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08)<br>Paper No(s)/Mail Date _____ | 6) <input type="checkbox"/> Other: _____  |

## **DETAILED ACTION**

1. This written action is responding to the amendment dated 08/07/2006.
2. Claims 34, 36, 43-4, 48, 51-53, 57-59, 62-64, and 66 have been amended, Claims 35 and 49 have been cancelled, and Claims 67-73 are newly added claims.
3. Claims 34-73 have been examined and rejected.
4. Claims 34-73 are pending
5. Rejections of Independent claims are provided with detailed citations from the prior arts.

## **Claim Objections**

6. Claims 37-39 and 63-65 are objected to because of the following informalities:
  - a. Claim 37 contains the limitation, "encrypted using the key", which appears to be a typographical error. The correct limitation should be "encrypted using the leaf key". For examining purpose, the claim is treated with the suggested correct term. Appropriate correction is required.
  - b. Claims 38 and 39 contain the limitation, "encrypted using the key", which appears to be a typographical error. The correct limitation should be "encrypted using the leaf key". For examining purpose, the claim is treated with the suggested correct term. Appropriate correction is required.

Art Unit: 2135

- c. Claim 63 on page 10 of the amended claim contains incomplete claim limitations. For examining purpose, the missing claim limitations are drawn from the previously submitted claims, dated on Feb. 24, 2006. Appropriate correction is required.
- d. Claim 64 on page 11 of the amended claim contains incomplete claim limitations at the preamble. For examining purpose, the missing claim limitations are drawn from the previously submitted claims, dated on Feb. 24, 2006. Appropriate correction is required.
- e. Claim 65 is objected for the deficiency inherited from Claim 64. Appropriate correction is required.

### **Claim Rejections - 35 USC § 103**

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

- 7. Claims 34, 36, 38-42, 44-48, 50, 52-56, 58-62, 64-65, 68, and 70-73 are rejected under 35 U.S.C. 103(a) as being unpatentable over Itkis (U.S. Patent 6,880,081) and further in view of Harada et al. (U.S. Patent 6,687,683) and Dondeti et al. (U.S. Patent 6,240,188).

a. Referring to Claims 34 and 44:

As per Claim 34, Itkis discloses an information processing device operable within a node of a hierarchical network of nodes having a hierarchical tree structure **[(Fig. 2)]**, said information processing device comprising:

to store one or more node keys, each node key being unique to one node of the network, and a leaf key, the leaf key being unique to the information processing device **[In a preferable implementation of the group assignments 20 as shown in FIG. 1, the group assignments 20 may be depicted as a tree in which each one of the plurality of authorized devices is represented by a leaf (lines 21-26, Col. 8). At level n, the leaf level, each group 100 is associated with a device 110 (lines 16-17, Col. 9 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2 (lines 41-44, Col. 9 and Fig. 2 from Itkis)]**;

to perform encryption processing **[It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all**

**groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 from Itkis)] to:**

**calculate a decryption key by decrypting a key block using at least one of the one or more node keys stored in the storage or the leaf key stored in the storage [Accompanying the content is a key block B (the key block can be assumed to include "media key" – e.g., the disc's serial number, etc. (lines 51-53, Col. 1 from Itkis). B can be computed (by the content providers, after examining the pirate devices) in such a way that all non-compromised devices can compute K from B (lines 56-58, Col. 1 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 and Fig. 2 from Itkis). Where K is a content encryption key or any other useful key (lines 7-8, Col. 10 from Itkis). Thus, each device 110 need only perform one decryption operation in order to obtain K. It is appreciated that a further, typically fixed number of decryption operations, as is well known in the art, may need to be performed in order to actually obtain protected content (lines 12-16, Col. 10 from Itkis)];**

encrypt the decryption key using the key of the information processing device [At level n, the leaf level, each group 100 is associated with a device 110 (lines 16-17, Col. 9 from Itkis). Where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 and Fig. 2 from Itkis)].

Itkis discloses the hardware component for performing the encryption, decryption, and storage processes [In a preferred embodiment of the present invention, an improved key distribution system is provided (lines 49-50, Col. 2 from Itkis). Each of the components of FIG. 1 is preferably implemented in a combination of software and computer hardware, as is well known in the art, and may include special purpose computer hardware, as is also well known in art, in order to increase efficiency of operation (lines 3-7, Col. 8 and Fig. 1). Individual components, described below, of the security element 120 may be implemented in hardware or in any suitable combination of hardware and software, as is well known in the art (lines 5-8, Col. 11 and Fig. 4)].

However, Itkis does not expressly disclose (1) storage and encryption processor within the processing devices for holding the key information, calculate decryption key, and executing

encryption/decryption processing, (2) store the encrypted decryption key in at least one of the storage or on a recording medium. However, Harada et al. disclose (1) the LSI component, which contains the encryption and decryption units for deriving key information and performing encryption/decryption processes with keys and the storage unit for holding the relevant key information [The disk key creation unit 1218 creates a 64-bit disk key including the information on the memory card ID that has been given from the memory card ID obtaining unit 1230. Here, a disk key is key data common to all kinds of memory card that is recording medium. The disk key encryption unit 1220 encrypts the disk key that has been created by the disk key creation unit 1218 using one of the plurality of master keys 1219 that have been stored in the disk key encryption unit 1220 in advance. The disk key encryption unit 1220 continues to encrypt the same disk key using a different master key 1219 to create the same number of encryption disk keys as that of the master keys 1219, and outputs the created encryption disk keys to the recording unit 1240 in the memory card writer 1200. The title key creation unit 1221 creates an appropriate 64-bit title key and gives the created title key to the title key encryption unit 1222. Here, the title key indicates key data that can be set for



each music content (lines 8-24, Col. 13 from Harada et al.). Meanwhile, the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 29-34, Col. 13 and unit 1200 in Fig. 2 and 3 from Harada et al.)), and (2) to store the encrypted decrypting key on a recording medium or in a storage area in said information processing device [The title key encryption unit 1222 encrypts the title key that has been created by the title key creation unit 1221 using the disk that has been created by the disk key creation unit 1218, and outputs the encrypted title key to the recording unit 1240. Meanwhile the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 25-34, Col. 13). Note that the recording unit 1240 records the audio data that has been transferred from the audio data encryption unit 1223 in an user accessible area in the memory card 1300 and the encrypted disk key and title key in a system area in the memory card 1300 that cannot be

accessed by the user (lines 39-43, Col. 13 and Fig. 2 from Harada et al.)). Dondeti et al. further disclose the leaf key being unique to each of said information processing devices and is unique in relation to a leaf key held by any other node within the hierarchical network of nodes [Each member 22 is assigned a binary ID and these IDs are used to define key associations for each member 22 (lines 30-31, Col. 3). Members are represented by the leaves of a binary key distribution tree 26. Each member 22 generates a unique secret key 28 for itself and each internal node key is computed as a function of the secret keys of its two children. All secret keys 28 are associated with their blinded versions 30, which are computed using a one-way function 32 (lines 48-53, Col. 3). Internal nodes are associated with secret keys (lines 2-3, Col. 4)], and the key used in encrypting the calculated decrypting key is the leaf key [Wherein, the first member uses the blinded keys received from the key association group and the first secret key to calculate an unblinded key of the first internal node (lines 48-51, Col. 2). All secret keys 28 are associated with their blinded versions 30, which are computed using a one-way function 32 (lines 52-53, Col. 3)]. Itkis, Harada et al., and Dondeti et al. are analogous art because they are from similar technology relating to the digital content information

security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis with Harada et al. and Dondeti et al. since one would have been motivated to (1) provide a production protection system that enables contents to be recorded on a recording medium loaded on a player for replaying contents and the like in order to more intensively protect contents for sale, and enables one of encryption algorithms for distributing contents via the Internet and for recording contents on the recording medium not to be influenced by the decryption of the other one (lines 53-60, Col. 1 from Harada et al.) and (2) to have a system for providing secure communication between many sends and many members (lines 15-17 from Dondeti et al.). Therefore, it would have been obvious to combine Itkis with Harada et al. and Dondeti et al. to obtain the invention as specified in claim 34.

As per Claim 44, it encompasses some limitations that are similar to those of Claim 34. Therefore, these limitations are rejected with the same rationale applied against Claim 34 above. In addition, Harada et al. disclose store the decryption key together with identification information, the identification information being usable to identify data decrypted using the decryption key [the memory card ID obtaining unit 1230 obtains the memory card

ID that is inherent information from the memory card 1300, and gives the obtained memory card ID to the disk key creation unit 1218. When receiving the recording allowance, the recording unit 1240 records that data that have been output from the disk key encryption unit 1220, the title key encryption unit 1222, and the audio data encryption unit 1223 on the memory card 1300 (line 67, Col. 12 and lines 1-7, Col. 13 from Harada et al.). Meanwhile, the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221 (lines 29-32, Col. 13 from Harada et al.)).

b. Referring to Claim 36:

As per Claim 3, Itkis, Harada et al., and Dondeti et al. the information processing device as claimed in claim 34, wherein the key block includes an encrypted renewal node key, and is further operable to decrypt the encrypted renewal node key to obtain the renewal node key using at least one of the node key stored in the storage or a leaf key belonging to a lower layer of the hierarchical network, and stored in the storage [(lines 51-53, Col. 1; lines 56-59, Col. 2; lines 56-60, Col. 2; lines 9-10 and 54-58, Col. 3; lines 49-56, Col. 9 from Itkis) and (lines 48-53, Col. 3; lines 2-3, Col. 4 from Dondeti et al.)], and to calculate

the decryption key using the obtained renewal node key **[(lines 55-59, Col. 1 and lines 7-18, Col. 10 from Itkis)]**. In addition, Harada et al. disclose the encryption processor as in claim 34.

c. Referring to Claims 38 and 52:

As per Claim 38, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the encryption processor is operable to store the decryption key encrypted using the leaf key unique to the information processing device, **[(lines 54-59, Col. 3 and lines 43-48, Col. 9 from Itkis) and (lines 48-53, Col. 3; lines 2-3, Col. 4 from Dondeti et al.)]**, the encrypted decryption key being stored together with identification information, the identification information being unique to the information processing device **[(lines 7-11, Col. 10 from Itkis)]**.

As per Claim 52, the rejection of Claim 48 is incorporated. In addition, Claim 52 encompasses limitations that are similar to those of Claim 38. Therefore, it is rejected with the same rationale applied against Claim 38 above.

d. Referring to Claims 39 and 53:

As per Claim 39, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the encryption processing is operable to store the decryption key

encrypted using the leaf key unique to the information processing device, the encrypted decryption key being stored together with identification information, the identification information identifying data decrypted using the decryption key **[(line 67, Col. 12 and lines 1-7 and 29-32 Col. 13 from Harada et al.) and (lines 48-53, Col. 3; lines 2-3, Col. 4 from Dondeti et al.)]**.

As per Claim 53, the rejection of Claim 48 is incorporated. In addition, Claim 53 encompasses limitations that are similar to those of Claim 39. Therefore, it is rejected with the same rationale applied against Claim 39 above.

e. Referring to Claims 40 and 54:

As per Claim 40, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the decryption key is usable to decrypt encrypted content data in the information processing device **[(lines 48-50, Col. 1 and lines 7-8, Col. 10 from Itkis)]**.

As per Claim 54, the rejection of Claim 48 is incorporated. In addition, Claim 54 encompasses limitations that are similar to those of Claim 40. Therefore, it is rejected with the same rationale applied against Claim 40 above.

f. Referring to Claims 41 and 55:

Art Unit: 2135

As per Claim 41, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the decryption key is stored on the recording medium and the decryption key is assigned to the recording medium, the decryption key being usable to decrypt encrypted data stored on the recording medium **[(lines 48-50, Col. 1, and lines 7-8, Col. 10 from Itkis) and (lines 25-34 and 39-43, Col. 13 from Harada et al.); where the title key is the (media) content key and is located in the recording medium]**.

As per Claim 55, the rejection of Claim 48 is incorporated. In addition, Claim 55 encompasses limitations that are similar to those of Claim 41. Therefore, it is rejected with the same rationale applied against Claim 41 above.

g. Referring to Claims 42 and 56:

As per Claim 42, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the decryption key is held in common by a plurality of the information processing devices, the decryption key being a master key usable to decrypt processing devices **(lines 48-50, Col. 1, lines 7-8, Col. 10, and lines 44-48, Col. 9 from Itkis); where K is always the key used for decrypting the content even it is encrypted by different versions of keys]**.

As per Claim 56, the rejection of Claim 48 is incorporated. In addition, Claim 56 encompasses limitations that are similar to those of Claim 42. Therefore, it is rejected with the same rationale applied against Claim 42 above.

h. Referring to Claims 45, 59, and 64:

As per Claim 45, Itkis discloses an information processing device operable within a node of a hierarchical network of nodes having a hierarchical tree structure **[(Fig. 2)]**, said information processing device comprising:

to store a node key and a leaf key, the leaf key being unique to the information processing device, and the node key being unique to each node of a hierarchical network of nodes having a hierarchical tree structure **[In a preferable implementation of the group assignments 20 as shown in FIG. 1, the group assignments 20 may be depicted as a tree in which each one of the plurality of authorized devices is represented by a leaf (lines 21-26, Col. 8). At level n, the leaf level, each group 100 is associated with a device 110 (lines 16-17, Col. 9 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2 (lines 41-44, Col. 9 and Fig. 2 from Itkis)]; and**



Art Unit: 2135

operable to perform decryption processing **[Thus each device 110 need only perform one decryption operation in order to obtain K. It is appreciated that a further, typically fixed number of decryption operation, as is well known in the art, may need to be performed in order to actually obtain protected content (lines 12-16, Col. 10 from Itkis)]** to:

detect whether an encrypted decryption key for decrypting encrypted data, and when the encrypted decryption key is detected, to calculate the decryption key by decrypting the encrypted decryption key **[performing no more than a predetermined number of decryption operations, the predetermined number being the same for all authorized devices, to obtain the content decryption key from an encrypted form thereof, the encrypted form being encrypted with a group key corresponding to a group of which the authorized device is a member (lines 54-59, Col. 3 from Itkis). A key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 670 are used, independently, to encrypt K (lines 43-46, Col. 9 from Itkis)], and**

when the encrypted decryption key is not detected, to calculate the decryption key by decrypting a key block using at least one of the one or more node keys stored in the storage or the leaf key

stored in the storage [Accompanying the content is a key block B (the key block can be assumed to include "media key" – e.g., the disc's serial number, etc. (lines 51-53, Col. 1 from Itkis). B can be computed (by the content providers, after examining the pirate devices) in such a way that all non-compromised devices can compute K from B (lines 56-58, Col. 1 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 and Fig. 2 from Itkis). Where K is a content encryption key or any other useful key (lines 7-8, Col. 10 from Itkis). Thus, each device 110 need only perform one decryption operation in order to obtain K. It is appreciated that a further, typically fixed number of decryption operations, as is well known in the art, may need to be performed in order to actually obtain protected content (lines 12-16, Col. 10 from Itkis)].

Itkis discloses the hardware component for performing the encryption, decryption, and storage processes [In a preferred embodiment of the present invention, an improved key distribution system is provided (lines 49-50, Col. 2 from Itkis).

Each of the components of FIG. 1 is preferably implemented in a combination of software and computer hardware, as is well known in the art, and may include special purpose computer hardware, as is also well known in art, in order to increase efficiency of operation (lines 3-7, Col. 8 and Fig. 1). Individual components, described below, of the security element 120 may be implemented in hardware or in any suitable combination of hardware and software, as is well known in the art (lines 5-8, Col. 11 and Fig. 4)]. Itkis does not expressly disclose the hardware containing: (1) storage and encryption processor within the processing devices for holding the key information, calculate decryption key, and executing encryption/decryption processing, (2) the encrypted decrypting key is stored on the recording medium or in the recording area in the information processing device. However, Harada et al. disclose an LSI component, which contains the encryption and decryption units for deriving key information and performing encryption/decryption processes with keys and a storage unit for holding the relevant key information [The disk key creation unit 1218 creates a 64-bit disk key including the information on the memory card ID that has been given from the memory card ID obtaining unit 1230. Here, a disk key is key data common to all kinds of memory card that is recording

medium. The disk key encryption unit 1220 encrypts the disk key that has been created by the disk key creation unit 1218 using one of the plurality of master keys 1219 that have been stored in the disk key encryption unit 1220 in advance. The disk key encryption unit 1220 continues to encrypt the same disk key using a different master key 1219 to create the same number of encryption disk keys as that of the master keys 1219, and outputs the created encryption disk keys to the recording unit 1240 in the memory card writer 1200. The title key creation unit 1221 creates an appropriate 64-bit title key and gives the created title key to the title key encryption unit 1222. Here, the title key indicates key data that can be set for each music content (lines 8-24, Col. 13 from Harada et al.). Meanwhile, the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 29-34, Col. 13 and unit 1200 in Fig. 2 and 3 from Harada et al.)), and the encrypted decrypting key stored on a recording medium or in a storage area in said information processing device [The title key encryption unit 1222 encrypts the title key that has been created by the title key creation unit 1221 using the disk

that has been created by the disk key creation unit 1218, and outputs the encrypted title key to the recording unit 1240. Meanwhile the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 25-34, Col. 13). Note that the recording unit 1240 records the audio data that has been transferred from the audio data encryption unit 1223 in an user accessible area in the memory card 1300 and the encrypted disk key and title key in a system area in the memory card 1300 that cannot be accessed by the user (lines 39-43, Col. 13 and Fig. 2 from Harada et al.)). Dondeti et al. further disclose the leaf key being unique to the information processing device and unique in relation to a leaf key held by any other node within the hierarchical network of nodes [Each member 22 is assigned a binary ID and these IDs are used to define key associations for each member 22 (lines 30-31, Col. 3). Members are represented by the leaves of a binary key distribution tree 26. Each member 22 generates a unique secret key 28 for itself and each internal node key is computed as a function of the secret keys of its two children. All secret keys 28 are associated

with their blinded versions 30, which are computed using a one-way function 32 (lines 48-53, Col. 3). Internal nodes are associated with secret keys (lines 2-3, Col. 4)]. Itkis, Harada et al., and Dondeti et al. are analogous art because they are from similar technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis with Harada et al. and Dondeti et al. since one would have been motivated to (1) provide a production protection system that enables contents to be recorded on a recording medium loaded on a player for replaying contents and the like in order to more intensively protect contents for sale, and enables one of encryption algorithms for distributing contents via the Internet and for recording contents on the recording medium not to be influenced by the decryption of the other one (lines 53-60, Col. 1 from Harada et al.) and (2) to have a system for providing secure communication between many sends and many members (lines 15-17 from Dondeti et al.). Therefore, it would have been obvious to combine Itkis with Harada et al. and Dondeti et al. to obtain the invention as specified in claim 45.

As per Claim 59, it is a method claim that encompasses limitations that are similar to those of the device Claim 45.

Therefore, it is rejected with the same rationale applied against Claim 45 above.

As per Claims 64, it is a recording medium containing computer program claim corresponding to the method claim 59. Thus, it is rejected with the same rationale applied against Claim 59 above. In addition, Harada et al. disclose the computer program executed on an information processing device **[The personal computer 1100 is a personal computer that includes a CPU, a memory, a hard disk and the like and executes a program (lines 43-45, Col. 7 from Harada et al.)]**.

i. Referring to Claims 46, 60, and 65:

As per Claim 46, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 45, wherein, when the decryption key is not detected the decryption processor is further operable to encrypt the calculated decryption key and to store the encrypted decryption key on at least one of the recording medium or the memory **[(lines 12-15, 25-29, and 39-43, Col. 13 from Harada et al.)]**.

As per Claim 60, the rejection of Claim 59 is incorporated. In addition, Claim 60 encompasses limitations that are similar to those of Claim 46. Therefore, it is rejected with the same

Art Unit: 2135

rationale applied against Claim 46 above. In addition, Itkis discloses said decrypting key calculated using at least one of the node key and the leaf key held in said storage means **[(lines 16-17 and 43-48, Col. 9 and Fig. 2 from Itkis)]**.

As per Claim 65, the rejection of Claim 64 is incorporated. In addition, Claim 65 is a recording medium claim corresponding to the method Claim 60. Therefore, it is rejected with the same rationale applied against Claim 60 above.

j. Referring to Claims 47 and 61:

As per Claim 47, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 45, wherein the decryption processor is further operable to decrypt the encrypted decryption key using at least one key unique to the information processing device when the encrypted decryption key is detected **[(lines 54-59, Col. 3 and lines 43-48, Col. 9 from Itkis)]**.

As per Claim 61, the rejection of Claim 59 is incorporated. In addition, Claim 61 encompasses limitations that are similar to those of Claim 47. Therefore, it is rejected with the same rationale applied against Claim 47 above.

k. Referring to Claims 48 and 62:



Art Unit: 2135

As per Claim 48, Itkis discloses an information processing method; comprising:

storing one or more node keys and a leaf key in an information processing device of one node of a hierarchical network of nodes having a hierarchical tree structure, each node key being unique to one node of the network, the leaf key being unique to the information processing device **[In a preferable implementation of the group assignments 20 as shown in FIG. 1, the group assignments 20 may be depicted as a tree in which each one of the plurality of authorized devices is represented by a leaf (lines 21-26, Col. 8). At level n, the leaf level, each group 100 is associated with a device 110 (lines 16-17, Col. 9 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2 (lines 41-44, Col. 9 and Fig. 2 from Itkis)];**

decrypting a key block using at least one of the stored node key and the stored leaf key **[Accompanying the content is a key block B (the key block can be assumed to include "media key" – e.g., the disc's serial number, etc. (lines 51-53, Col. 1 from Itkis). B can be computed (by the content providers, after examining the pirate devices) in such a way that all non-compromised devices can compute K from B (lines 56-58,**

Col. 1 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 and Fig. 2 from Itkis). Where K is a content encryption key or any other useful key (lines 7-8, Col. 10 from Itkis). Thus, each device 110 need only perform one decryption operation in order to obtain K. It is appreciated that a further, typically fixed number of decryption operations, as is well known in the art, may need to be performed in order to actually obtain protected content (lines 12-16, Col. 10 from Itkis)];

calculating a decryption key usable to decrypt encrypted data stored on at least one of the information processing device or on a recording medium [K may be typically be obtained from B in the present invention by a legitimate device in a single decryption operation (lines 58-60, Col. 2). Where K is a content encryption key or any other useful key (lines 7-8, Col. 10)]; and

encrypting the decryption key using the key of the information processing device [It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key

**distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2. At any point, the keys of all groups 100 in the authorized set 60 are used, independently, to encrypt K (lines 41-46, Col. 9 from Itkis)].**

Itkis does not expressly disclose storing the encrypted decryption key on at least one of the information processing device or on the recording medium and other limitations of the claim. However, Harada et al. disclose the encrypted decryption key, used for decrypting the content information, is stored on the recording medium [The title key encryption unit 1222 encrypts the title key that has been created by the title key creation unit 1221 using the disk that has been created by the disk key creation unit 1218, and outputs the encrypted title key to the recording unit 1240. Meanwhile the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 25-34, Col. 13). Note that the recording unit 1240 records the audio data that has been transferred from the audio data encryption unit 1223 in an user accessible area in the memory card 1300 and the encrypted disk key and title key in a system area in the memory card 1300 that cannot be

**accessed by the user (lines 39-43, Col. 13 and Fig. 2)].**

Dondeti et al. further disclose the leaf key being unique to the information processing device such that each leaf key of each information processing device of the network is unique with respect to a leaf key of any other information processing device of the network **[Each member 22 is assigned a binary ID and these IDs are used to define key associations for each member 22 (lines 30-31, Col. 3). Members are represented by the leaves of a binary key distribution tree 26. Each member 22 generates a unique secret key 28 for itself and each internal node key is computed as a function of the secret keys of its two children. All secret keys 28 are associated with their blinded versions 30, which are computed using a one-way function 32 (lines 48-53, Col. 3). Internal nodes are associated with secret keys (lines 2-3, Col. 4)], and the key used in encrypting the calculated decrypting key is the leaf key [Wherein, the first member uses the blinded keys received from the key association group and the first secret key to calculate an unblinded key of the first internal node (lines 48-51, Col. 2). All secret keys 28 are associated with their blinded versions 30, which are computed using a one-way function 32 (lines 52-53, Col. 3)]. Itkis, Harada et al., and Dondeti et al. are analogous art because they are from similar**

Art Unit: 2135

technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis with Harada et al. and Dondeti et al. since one would have been motivated to (1) provide a production protection system that enables contents to be recorded on a recording medium loaded on a player for replaying contents and the like in order to more intensively protect contents for sale, and enables one of encryption algorithms for distributing contents via the Internet and for recording contents on the recording medium not to be influenced by the decryption of the other one (lines 53-60, Col. 1 from Harada et al.) and (2) to have a system for providing secure communication between many sends and many members (lines 15-17 from Dondeti et al.). Therefore, it would have been obvious to combine Itkis with Harada et al. and Dondeti et al. to obtain the invention as specified in claim 48.

As per Claim 62, it is a recording medium containing computer program claim corresponding to the method claim 48. Thus, it is rejected with the same rationale applied against Claim 48 above. In addition, Harada et al. disclose the computer program executed on an information processing device **[The personal computer 1100 is a personal computer that includes a CPU, a memory,**

**a hard disk and the like and executes a program (lines 43-45, Col. 7 from Harada et al.).**

**l. Referring to Claim 50:**

As per Claim 50, Itkis, Harada et al., and Dondeti et al. disclose the information processing method as claimed in claim 48, wherein the key block includes a renewal node key, the renewal node key being encrypted using at least one of the stored node key for the node or a leaf key belonging to a lower layer of the hierarchical network **[(lines 51-53, Col. 1; lines 56-59, Col. 2, lines 9-10; 54-58, Col. 3; lines 49-56, Col. 9 from Itkis)]**, and the decryption key is encrypted using the renewal node key **[(lines 58-60, Col. 2 from Itkis)]**, wherein the step of decrypting the key block includes decrypting the renewal node key using at least one of the stored node key and the stored leaf key, and the calculating step includes using the decrypted renewal node key to calculate the decryption key **[(lines 55-62, Col. 1 and lines 7-18, Col. 10 from Itkis)]**.

**m. Referring to Claims 58:**

As per Claim 58, it encompasses some limitations that are similar to those of Claim 48. Therefore, these limitations are rejected with the same rationale applied against Claim 48 above. In addition, Harada et al. disclose storing the calculated decrypting key in the information processing device together with

identification information, the identification information being usable to identify data decrypted using said decrypting key [the memory card ID obtaining unit 1230 obtains the memory card ID that is inherent information from the memory card 1300, and gives the obtained memory card ID to the disk key creation unit 1218. When receiving the recording allowance, the recording unit 1240 records that data that have been output from the disk key encryption unit 1220, the title key encryption unit 1222, and the audio data encryption unit 1223 on the memory card 1300 (line 67, Col. 12 and lines 1-7, Col. 13 from Harada et al.). Meanwhile, the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221 (lines 29-32, Col. 13 from Harada et al.)].

n. Referring to Claim 68:

As per Claim 68, Itkis, Harada et al., and Dondeti et al. disclose the recording medium as claimed in claim 62, wherein the method further comprises accessing the stored encrypted decryption key, recovering the decryption key by decrypting the encrypted decryption key using the leaf key, and decrypting content information stored on at least one of the recording medium or the storage using the recovered decryption key [(lines 51-53 and 56-

**58, Col. 1; lines 12-16, Col. 10 from Itkis) and (lines 25-34 and 39-43, Col. 13; lines 37-48, Col. 16 from Harada et al.) and (lines 48-51, Col. 2 and lines 52-53, Col. 3 from Dondeti et al.)].**

**o. Referring to Claim 70:**

As per Claim 70, Itkis, Harada et al., and Dondeti et al. disclose the information processing apparatus as claimed in claim 34, wherein the encryption processor is further operable to access the stored encrypted decryption key, recover the decryption key by decrypting the encrypted decryption key using the leaf key, and decrypt content information stored on at least one of the recording medium or the storage using the recovered decryption key **[(lines 51-53 and 56-58, Col. 1; lines 12-16, Col. 10 from Itkis) and (lines 25-34 and 39-43, Col. 13; lines 37-48, Col. 16 from Harada et al.) and (lines 48-51, Col. 2 and lines 52-53, Col. 3 from Dondeti et al.)].**

**p. Referring to Claim 71:**

As per Claim 71, Itkis, Harada et al., and Dondeti et al. disclose the information processing apparatus as claimed in claim 70, wherein the decryption key includes a media key **[(lines 25-34 and 39-43, Col. 13 from Harada) and (lines 51-52, Col. 1 from Itkis)].**

**q. Referring to Claim 72:**



As per Claim 72, Itkis, Harada et al., and Dondeti et al. disclose the information processing method as claimed in claim 48, further comprising accessing the stored encrypted decryption key, recover the decryption key by decrypting the encrypted decryption key using the leaf key, and decrypt content information stored on at least one of the recording medium or the storage using the recovered decryption key **[(lines 51-53 and 56-58, Col. 1; lines 12-16, Col. 10 from Itkis) and (lines 25-34 and 39-43, Col. 13; lines 37-48, Col. 16 from Harada et al.) and (lines 48-51, Col. 2 and lines 52-53, Col. 3 from Dondeti et al.)]**.

r. Referring to Claim 73:

As per Claim 73, Itkis, Harada et al., and Dondeti et al. disclose the information processing method as claimed in claim 72, wherein the decryption key includes a media key **[(lines 25-34 and 39-43, Col. 13 from Harada) and (lines 51-52, Col. 1 from Itkis)]**.

8. Claims 66-67 and 69 are rejected under 35 U.S.C. 103(a) as being unpatentable over Harada et al. (U.S. Patent 6,687,683) in view of Itkis (U.S. Patent 6,880,081) and Dondeti et al. (U.S. Patent 6,240,188).

a. Referring to Claim 66:

As per Claim 66, Harada et al. disclose a recording medium having encrypted information recorded thereon including at least

one of audio information, video information or human language text information in encrypted form [The music content replay/recording system 1000 is a system in which a music content that has been received via a communication line 1001 is replayed using a personal computer 1100 and the music content is recorded on a memory card 1300 (lines 26-30, Col 7 from Harada et al.)], the encrypted information being decryptable only by any one of a plurality of information processing devices using a decryption key, the recording medium having the decryption key recorded thereon in encrypted form [The title key encryption unit 1222 encrypts the title key that has been created by the title key creation unit 1221 using the disk that has been created by the disk key creation unit 1218, and outputs the encrypted title key to the recording unit 1240. Meanwhile the audio data encryption unit 1223 re-encrypts the C2 content 40 that has been output from the C2 content decryption unit 1217 using the title key that has been created by the title key creation unit 1221, and outputs the re-encrypted C2 content 40 to the recording unit 1240 (lines 25-34, Col. 13). Note that the recording unit 1240 records the audio data that has been transferred from the audio data encryption unit 1223 in an user accessible area in the memory card 1300 and the encrypted disk key and title key in

a system area in the memory card 1300 that cannot be accessed by the user (lines 39-43, Col. 13 and Fig. 2 from Harada et al.). The encrypted C1 content 130 is data that is created by encrypting a plaintext, the C1 content 30 using the C1 key 21 (refer to FIG. 4), and the data length can change according to the content. The C1 key 21 is 40-bit key data, and the encryption using the C1 key 21 is performed in a block cipher system. For instance, a DES algorithm is used. The encrypted C2 content 140 is data that is created by encrypting a plaintext, the C2 content 40 using the C2 key 25 (refer to FIG. 4), and the data length can change according to the content. The C2 key 25 is 56-bit key data, and the encryption using the C2 key 25 is performed in a block cipher system. For instance, a DES algorithm is used (lines 37-48, Col. 16 from Harada et al.)). Harada et al. do not expressly disclose the remaining limitation of the claim. However, Itkis and Dondeti et al. disclose the encrypted decryption key having been encrypted using a leaf key unique to one information processing device, the encrypted decryption key being stored as a key storage table together with identification for the one information processing device [[In a preferable implementation of the group assignments 20 as shown in FIG. 1, the group assignments 20 may be depicted as a tree in which each one

of the plurality of authorized devices is represented by a leaf (lines 21-26, Col. 8). At level n, the leaf level, each group 100 is associated with a device 110 (lines 16-17, Col. 9 from Itkis). It will be appreciated that the system of FIG. 1 is particularly useful as a solution of the key distribution problem in a case where a key is assigned to each of the groups 100 of FIG. 2 (lines 41-44, Col. 9 and Fig. 2 from Itkis)] and [Each member 22 is assigned a binary ID and these IDs are used to define key associations for each member 22 (lines 30-31, Col. 3). Members are represented by the leaves of a binary key distribution tree 26. Each member 22 generates a unique secret key 28 for itself and each internal node key is computed as a function of the secret keys of its two children. All secret keys 28 are associated with their blinded versions 30, which are computed using a one-way function 32 (lines 48-53, Col. 3). Internal nodes are associated with secret keys (lines 2-3, Col. 4)]], and Itkis further discloses a key storage table together with identification for the one information processing device [1. Where K is a content encryption key or any other useful key, for example, device 110 can easily determine, based on group membership of the device 110 and, preferably, group identification accompanying each encryption of K in a key block B (lines 7-11, Col. 10 from

**Itkis]**. Harada et al., Itkis, and Dondeti et al. are analogous art because they are from similar technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Harada et al. with Itkis and Dondeti et al. since one would have been motivated to (1) provide improved apparatus and methods for content access control and improved key distribution system (lines 47-50, Col. 2 from Itkis) and (2) to have a system for providing secure communication between many sends and many members (lines 15-17 from Dondeti et al.). Therefore, it would have been obvious to combine Harada et al. with Itkis and Dondeti et al. to obtain the invention as specified in claim 66.

b. Referring to Claim 67:

As per Claim 67, Harada et al., Itkis, and Dondeti et al. disclose the recording medium as claimed in claim 66, wherein the recording medium is removably insertable into any information processing device of the plurality of information processing devices through an opening in an exterior housing of such information processing device, and is recordable to store the encrypted decryption key when the recording medium is inserted into the one information processing device [(lines 25-34 and 39-

**43, Col. 13; lines 37-48, Col. 16; Fig. 1 and 2 from Harada et al.)).**

c. Referring to Claim 69:

As per Claim 69, Harada et al., Itkis, and Dondeti et al. disclose the recording medium as claimed in claim 67, wherein the decryption key includes a media key **[(lines 25-34 and 39-43, Col. 13 from Harada) and (lines 51-52, Col. 1 from Itkis)].**

9. Claims 37, 43, 51, 57, and 63 are rejected under 35 U.S.C. 103(a) as being unpatentable over Itkis (U.S. Patent 6,880,081), Harada et al. (U.S. Patent 6,687,683), and Dondeti et al. (U.S. Patent 6,240,188), and further in view of Lotspiech et al. (U.S. Patent 6,118,873).

a. Referring to Claim 43:

As per Claim 43, it encompasses some limitations that are similar to those of Claim 34. Therefore, these limitations are rejected with the same rationale applied against Claim 34 above. Itkis, Harada et al., and Dondeti et al. do not expressly disclose the decryption key is with a generation number representing renewal information for the decrypting key. However, Lotspiech et al. disclose the renewal generation number is associated with the number of the time the decryption key has been renewed **[the renewal generation number refers to the number of times the keys of a device have been renewed (lines 21-23, Col. 6 from**

**Lotspiech et al.]).** Itkis, Harada et al., Dondeti et al., and Lotspiech et al. are analogous art because they are from similar technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis, Harada et al., Dondeti et al. with Lotspiech et al. since one would have been motivated to prevent the unauthorized viewing and/or copying (line 13, Col. 1 from Lotspiech et al.). Therefore, it would have been obvious to combine Itkis, Harada et al., and Dondeti et al. with Lotspiech et al. to obtain the invention as specified in claim 43.

b. Referring to Claims 37 and 51:

As per Claim 37, Itkis, Harada et al., and Dondeti et al. disclose the information processing device as claimed in claim 34, wherein the encryption processor is operable to store the decryption key encrypted using the leaf key unique to the information processing device **[(lines 54-59, Col. 3 and lines 43-48, Col. 9 from Itkis) and (lines 48-53, Col. 3; lines 2-3, Col. 4 from Dondeti et al.)]**. Itkis, Harada et al., and Dondeti et al. do not expressly disclose the encrypted decryption key being stored together with a generation number, the generation number representing renewal information for decryption key. However, Lotspiech et al. disclose the renewal generation number is associated with the number of

the time the decryption key has been renewed [(lines 21-23, Col. 6 from Lotspiech et al.)]. Itkis, Harada et al. and Lotspiech et al. are analogous art because they are from similar technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis and Harada et al. with Lotspiech et al. since one would have been motivated to prevent the unauthorized viewing and/or copying (line 13, Col. 1 from Lotspiech et al.). Therefore, it would have been obvious to combine Itkis and Harada et al. with Lotspiech et al. to obtain the invention as specified in claim 37.

As per Claim 51, the rejection of Claim 48 is incorporated. In addition, Claim 51 encompasses limitations that are similar to those of Claim 37. Therefore, it is rejected with the same rationale applied against Claim 37 above.

c. Referring to Claims 57 and 63:

As per Claim 57, it encompasses some limitations that are similar to those of Claim 48. Therefore, these limitations are rejected with the same rationale applied against Claim 48 above. Itkis, Harada et al., and Dondeti et al. do not expressly disclose the calculated decrypting key in the information processing device is with a generation number, where the generation number



represents renewal information for the decrypting key. However, Lotspiech et al. disclose the renewal generation number is associated with the number of the time the decryption key has been renewed **[the renewal generation number refers to the number of times the keys of a device have been renewed (lines 21-23, Col. 6 from Lotspiech et al.)]**. Itkis, Harada et al. and Lotspiech et al. are analogous art because they are from similar technology relating to the digital content information security and protection. It would have been obvious to one of ordinary skill in the art at the time of invention was made to modify Itkis and Harada et al. with Lotspiech et al. since one would have been motivated to prevent the unauthorized viewing and/or copying (line 13, Col. 1 from Lotspiech et al.). Therefore, it would have been obvious to combine Itkis and Harada et al. with Lotspiech et al. to obtain the invention as specified in claim 57.

As per Claim 63, it is a recording medium containing computer program claim corresponding to the method claim 57. Thus, it is rejected with the same rationale applied against Claim 57 above. In addition, Harada et al. disclose the computer program executed on an information processing device **[The personal computer 1100 is a personal computer that includes a CPU, a memory,**

**a hard disk and the like and executes a program (lines 43-45,  
Col. 7 from Harada et al.)].**

### **Response to Arguments**

10. Applicant's amendment, filed on Aug. 07, 2006, has Claims 34, 36, 43-4, 48, 51-53, 57-59, 62-64, and 66 amended, Claims 35 and 49 cancelled, and Claims 67-73 newly added. Among these amended claims, Claims 34, 43-45, 48, 57-59, 62-64, and 66 are the independent claims that have been modified. This necessitates the new grounds of rejection. Please refer to the rejections above.

11. The new limitation on the leaf key of the information processing device being the unique to other leaf key held by other node in the hierarchical tree network has been added to the amended independent Claims 34, 43-45, 48, 57-59, 62-64.
12. The prior art by Dondeti et al. (U.S. Patent 6,240,188), which has been previously disclosed in the Form 892, is used, in combination, with other previously cited references for rejecting the amended claims with the above-mentioned limitation (lines 30-31 and 48-53 Col. 3; lines 2-3, Col. 4 of Dondeti et al.). Please refer to the rejections above for details.

### Conclusion

13. Accordingly, **THIS ACTION IS MADE FINAL**. See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire **THREE MONTHS** from the mailing date of this action. In the event a first reply is filed within **TWO MONTHS** of the mailing date of this final action and the advisory action is not mailed until after the end of the **THREE-MONTH** shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than **SIX MONTHS** from the date of this final action.

- a. Dondeti et al. (U.S. Patent 6,264,435) disclose a logical tree structure and method for managing membership in a multicast group to provide scalability and security from internal attacks. The structure defines key groups and subgroups, with each subgroup having a subgroup manager. Dual encryption allows the sender of the multicast data to manage distribution of a first set of encryption keys whereas the individual subgroup managers manage the distribution of a second set of encryption keys. The two key sets allow the sender to delegate much of the group management responsibilities without compromising security because a key from each set is required to access the multicast data. Security is further maintained via a method in which subgroup managers can be either member subgroup managers or participant subgroup managers. Access to both keys is provided to member subgroup managers whereas access to only one key is provided to participant subgroup managers. Nodes can be added without the need to generate a new encryption key at the top level which provides improved scalability.
- b. Ober (U.S. Patent 6,959,086) disclose a key management scheme for managing encryption keys in a cryptographic co-processor includes the first step of selecting a key from one of a symmetrical key type and an asymmetrical key type. Then, the key bit length is

Art Unit: 2135


selected. The key is then generated and, lastly, the key is represented in either an external form or an internal form.

14. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Yin-Chen Shaw whose telephone number is 571-272-8593. The examiner can normally be reached on 8:15 to 4:15 M-F. If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Kim Yen Vu can be reached on 571-272-3859. The fax phone number for the organization where this application or proceeding is assigned is 703-872-9306.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

YCS

Oct. 27, 2006



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